

Techniques for preparing eigenstates of fermionic Hamiltonians

Mária Kieferová

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University of Technology Sydney

- Antisymmetrizing the wavefunction

[Berry et al., Improved Techniques for Preparing Eigenstates of Fermionic Hamiltonians, NPJ QI 2018]

- Simulating time-dependent Hamiltonians

[Kieferová et al., Simulating the dynamics of time-dependent Hamiltonians with a truncated Dyson series, PRA 2019]

our recent review:

[Cao et al., Quantum Chemistry in the Age of Quantum Computing, Chem. Rev. 2019]

Quantum Computing for Chemistry and Material Science

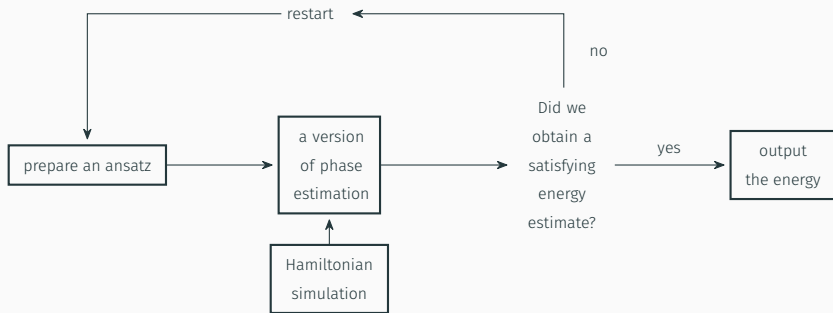


Algorithms for fault-tolerant quantum computers - beyond NISQ.

[credit: PharmaFactz, 123RF.com, Safety+Health Magazine]

Eigenenergy Estimation

Goal: compute the ground state energy of a molecule



[Abrams, Lloyd PRL (1997), Aspuru-Guzik et al. Science (2005)]

Antisymmetrization

First Quantization

State of m electrons, N sites/orbitals.

Fermions in the 1st quantization - symmetry is captured by the state.

- classically intractable
- typically better scaling than 2nd quantization for quantum algorithms

Preparation of Fermionic States

Input: ordered list of occupied orbitals

Output: completely antisymmetric state

$$|r_1 \cdots r_m\rangle \mapsto \sum_{\sigma \in S_m} (-1)^{\pi(\sigma)} |\sigma(r_1, \dots, r_m)\rangle$$

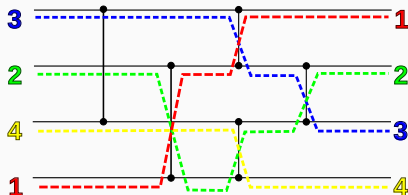
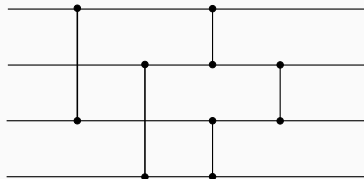
where $\pi(\sigma)$ is the parity of the permutation σ , and $r_p < r_{p+1}$.

NEW: Two approaches based on a classical (Fisher-Yates) shuffle and on reversed sort.

[Abrams, Lloyd PRL (1997)]

Sorting Network

Data-oblivious and reversible sort

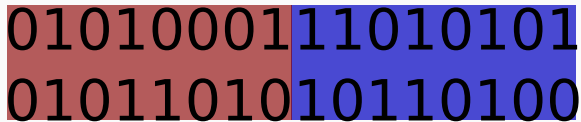


[Batcher 1968, Knuth 1968, source:wikipedia]

$$\begin{array}{c} A \\ B \end{array} \begin{array}{c} \bullet \\ \bullet \end{array} \begin{array}{c} \text{---} \\ \text{---} \end{array} = \begin{array}{c} A \\ B \end{array} \begin{array}{c} \diagup \text{---} \bullet \text{---} \times \\ \diagdown \text{---} \bullet \text{---} \times \end{array} \begin{array}{c} \min(A, B) \\ \max(A, B) \end{array} \\
 \begin{array}{c} |0\rangle \end{array} \text{---} \boxed{A > B} \begin{array}{c} \bullet \\ \bullet \end{array}$$

Bitonic sort requires $\mathcal{O}(m \log^2 m)$ comparators.

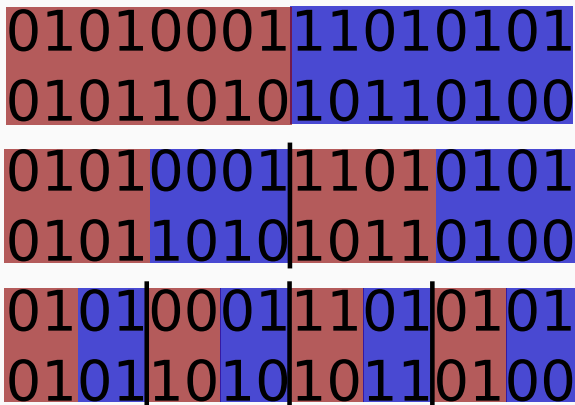
Coherent Comparison - Divide and Conquer



01010001	11010101
01011010	10110100

Compare each part separately and then merge the outcomes for the final comparison.

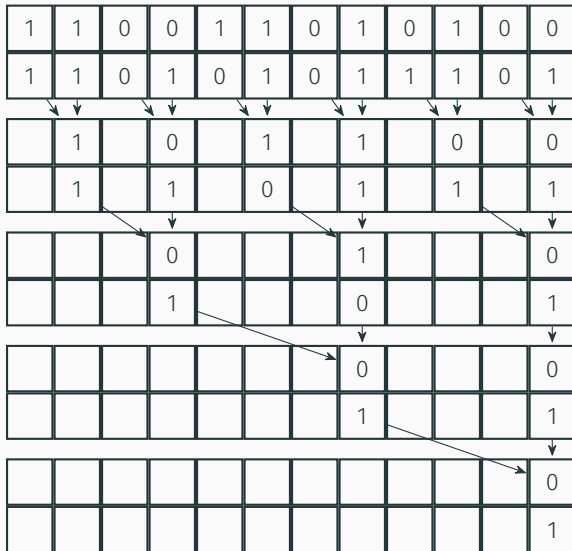
Coherent Comparison - Divide and Conquer



Only ever compare pairs of bits.

Gate count $\mathcal{O}(\log N)$, depth $\mathcal{O}(\log \log N)$.

How Does the Comparison work?



Reverse the algorithm to uncompute

Antisymmetrization via a Quantum Sort

$ 0\rangle 0\rangle$	→	$ 0\rangle 0\rangle$	→	$ 0\rangle 0\rangle \text{no swap} \rangle$	→	$0\rangle 0\rangle \text{no swap} \rangle$
	→	$ 0\rangle 1\rangle$	→	$ 0\rangle 1\rangle \text{no swap} \rangle$	→	$ 0\rangle 1\rangle \text{no swap} \rangle$
	→	$ 0\rangle 2\rangle$	→	$ 0\rangle 2\rangle \text{no swap} \rangle$	→	$ 0\rangle 2\rangle \text{no swap} \rangle$
	→	$ 0\rangle 3\rangle$	→	$ 0\rangle 3\rangle \text{no swap} \rangle$	→	$ 0\rangle 3\rangle \text{no swap} \rangle$
	→	$ 1\rangle 0\rangle$	→	$ 0\rangle 1\rangle \text{swap} \rangle$	→	$ 0\rangle 1\rangle \text{swap} \rangle$
	→	$ 1\rangle 1\rangle$	→	$ 1\rangle 1\rangle \text{no swap} \rangle$	→	$1\rangle 1\rangle \text{no swap} \rangle$
	→	$ 1\rangle 2\rangle$	→	$ 1\rangle 2\rangle \text{no swap} \rangle$	→	$ 1\rangle 2\rangle \text{no swap} \rangle$
	→	$ 1\rangle 3\rangle$	→	$ 1\rangle 3\rangle \text{no swap} \rangle$	→	$ 1\rangle 3\rangle \text{no swap} \rangle$
	→	$ 2\rangle 0\rangle$	→	$ 0\rangle 2\rangle \text{swap} \rangle$	→	$ 0\rangle 2\rangle \text{swap} \rangle$
	→	$ 2\rangle 1\rangle$	→	$ 1\rangle 2\rangle \text{swap} \rangle$	→	$ 1\rangle 2\rangle \text{swap} \rangle$
	→	$ 2\rangle 2\rangle$	→	$ 2\rangle 2\rangle \text{no swap} \rangle$	→	$2\rangle 2\rangle \text{no swap} \rangle$
	→	$ 2\rangle 3\rangle$	→	$ 2\rangle 3\rangle \text{no swap} \rangle$	→	$ 2\rangle 3\rangle \text{no swap} \rangle$
	→	$ 3\rangle 0\rangle$	→	$ 0\rangle 3\rangle \text{swap} \rangle$	→	$ 0\rangle 3\rangle \text{swap} \rangle$
	→	$ 3\rangle 1\rangle$	→	$ 1\rangle 3\rangle \text{swap} \rangle$	→	$ 1\rangle 3\rangle \text{swap} \rangle$
	→	$ 3\rangle 2\rangle$	→	$ 2\rangle 3\rangle \text{swap} \rangle$	→	$ 2\rangle 3\rangle \text{swap} \rangle$
	→	$ 3\rangle 3\rangle$	→	$ 3\rangle 3\rangle \text{no swap} \rangle$	→	$3\rangle 3\rangle \text{no swap} \rangle$

Create a uniform
superposition

Antisymmetrization via a Quantum Sort

$ 0\rangle 0\rangle$	→	$ 0\rangle 0\rangle$	→	$ 0\rangle 0\rangle$ no swap⟩	→	$0\rangle 0\rangle$ no swap⟩
	→	$ 0\rangle 1\rangle$	→	$ 0\rangle 1\rangle$ no swap⟩	→	$ 0\rangle 1\rangle$ no swap⟩
	→	$ 0\rangle 2\rangle$	→	$ 0\rangle 2\rangle$ no swap⟩	→	$ 0\rangle 2\rangle$ no swap⟩
	→	$ 0\rangle 3\rangle$	→	$ 0\rangle 3\rangle$ no swap⟩	→	$ 0\rangle 3\rangle$ no swap⟩
	→	$ 1\rangle 0\rangle$	→	$ 0\rangle 1\rangle$ swap ⟩	→	$ 0\rangle 1\rangle$ swap ⟩
	→	$ 1\rangle 1\rangle$	→	$ 1\rangle 1\rangle$ no swap⟩	→	$1\rangle 1\rangle$ no swap⟩
	→	$ 1\rangle 2\rangle$	→	$ 1\rangle 2\rangle$ no swap⟩	→	$ 1\rangle 2\rangle$ no swap⟩
	→	$ 1\rangle 3\rangle$	→	$ 1\rangle 3\rangle$ no swap⟩	→	$ 1\rangle 3\rangle$ no swap⟩
	→	$ 2\rangle 0\rangle$	→	$ 0\rangle 2\rangle$ swap ⟩	→	$ 0\rangle 2\rangle$ swap ⟩
	→	$ 2\rangle 1\rangle$	→	$ 1\rangle 2\rangle$ swap ⟩	→	$ 1\rangle 2\rangle$ swap ⟩
	→	$ 2\rangle 2\rangle$	→	$ 2\rangle 2\rangle$ no swap⟩	→	$2\rangle 2\rangle$ no swap⟩
	→	$ 2\rangle 3\rangle$	→	$ 2\rangle 3\rangle$ no swap⟩	→	$ 2\rangle 3\rangle$ no swap⟩
	→	$ 3\rangle 0\rangle$	→	$ 0\rangle 3\rangle$ swap ⟩	→	$ 0\rangle 3\rangle$ swap ⟩
	→	$ 3\rangle 1\rangle$	→	$ 1\rangle 3\rangle$ swap ⟩	→	$ 1\rangle 3\rangle$ swap ⟩
	→	$ 3\rangle 2\rangle$	→	$ 2\rangle 3\rangle$ swap ⟩	→	$ 2\rangle 3\rangle$ swap ⟩
	→	$ 3\rangle 3\rangle$	→	$ 3\rangle 3\rangle$ no swap⟩	→	$3\rangle 3\rangle$ no swap⟩

Sort registers and
record swaps

Antisymmetrization via a Quantum Sort

$ 0\rangle 0\rangle$	→	$ 0\rangle 0\rangle$	→	$ 0\rangle 0\rangle$ no swap	→	$0\rangle 0\rangle$ no swap
	→	$ 0\rangle 1\rangle$	→	$ 0\rangle 1\rangle$ no swap	→	$ 0\rangle 1\rangle$ no swap
	→	$ 0\rangle 2\rangle$	→	$ 0\rangle 2\rangle$ no swap	→	$ 0\rangle 2\rangle$ no swap
	→	$ 0\rangle 3\rangle$	→	$ 0\rangle 3\rangle$ no swap	→	$ 0\rangle 3\rangle$ no swap
	→	$ 1\rangle 0\rangle$	→	$ 0\rangle 1\rangle$ swap	→	$ 0\rangle 1\rangle$ swap
	→	$ 1\rangle 1\rangle$	→	$ 1\rangle 1\rangle$ no swap	→	$1\rangle 1\rangle$ no swap
	→	$ 1\rangle 2\rangle$	→	$ 1\rangle 2\rangle$ no swap	→	$ 1\rangle 2\rangle$ no swap
	→	$ 1\rangle 3\rangle$	→	$ 1\rangle 3\rangle$ no swap	→	$ 1\rangle 3\rangle$ no swap
	→	$ 2\rangle 0\rangle$	→	$ 0\rangle 2\rangle$ swap	→	$ 0\rangle 2\rangle$ swap
	→	$ 2\rangle 1\rangle$	→	$ 1\rangle 2\rangle$ swap	→	$ 1\rangle 2\rangle$ swap
	→	$ 2\rangle 2\rangle$	→	$ 2\rangle 2\rangle$ no swap	→	$2\rangle 2\rangle$ no swap
	→	$ 2\rangle 3\rangle$	→	$ 2\rangle 3\rangle$ no swap	→	$ 2\rangle 3\rangle$ no swap
	→	$ 3\rangle 0\rangle$	→	$ 0\rangle 3\rangle$ swap	→	$ 0\rangle 3\rangle$ swap
	→	$ 3\rangle 1\rangle$	→	$ 1\rangle 3\rangle$ swap	→	$ 1\rangle 3\rangle$ swap
	→	$ 3\rangle 2\rangle$	→	$ 2\rangle 3\rangle$ swap	→	$ 2\rangle 3\rangle$ swap
	→	$ 3\rangle 3\rangle$	→	$ 3\rangle 3\rangle$ no swap	→	$3\rangle 3\rangle$ no swap

Flag and delete
repetitions

Antisymmetrization via a Quantum Sort

	→	$ 0\rangle 0\rangle$	→	$ 0\rangle 0\rangle \text{no swap} \rangle$	→	$0\rangle 0\rangle \text{no swap} \rangle$
	→	$ 0\rangle 1\rangle$	→	$ 0\rangle 1\rangle \text{no swap} \rangle$	→	$ 0\rangle 1\rangle \text{no swap} \rangle$
	→	$ 0\rangle 2\rangle$	→	$ 0\rangle 2\rangle \text{no swap} \rangle$	→	$ 0\rangle 2\rangle \text{no swap} \rangle$
	→	$ 0\rangle 3\rangle$	→	$ 0\rangle 3\rangle \text{no swap} \rangle$	→	$ 0\rangle 3\rangle \text{no swap} \rangle$
	→	$ 1\rangle 0\rangle$	→	$ 0\rangle 1\rangle \text{swap} \rangle$	→	$ 0\rangle 1\rangle \text{swap} \rangle$
	→	$ 1\rangle 1\rangle$	→	$ 1\rangle 1\rangle \text{no swap} \rangle$	→	$1\rangle 1\rangle \text{no swap} \rangle$
	→	$ 1\rangle 2\rangle$	→	$ 1\rangle 2\rangle \text{no swap} \rangle$	→	$ 1\rangle 2\rangle \text{no swap} \rangle$
	→	$ 1\rangle 3\rangle$	→	$ 1\rangle 3\rangle \text{no swap} \rangle$	→	$ 1\rangle 3\rangle \text{no swap} \rangle$
$ 0\rangle 0\rangle$	→	$ 2\rangle 0\rangle$	→	$ 0\rangle 2\rangle \text{swap} \rangle$	→	$ 0\rangle 2\rangle \text{swap} \rangle$
	→	$ 2\rangle 1\rangle$	→	$ 1\rangle 2\rangle \text{swap} \rangle$	→	$ 1\rangle 2\rangle \text{swap} \rangle$
	→	$ 2\rangle 2\rangle$	→	$ 2\rangle 2\rangle \text{no swap} \rangle$	→	$2\rangle 2\rangle \text{no swap} \rangle$
	→	$ 2\rangle 3\rangle$	→	$ 2\rangle 3\rangle \text{no swap} \rangle$	→	$ 2\rangle 3\rangle \text{swap} \rangle$
	→	$ 3\rangle 0\rangle$	→	$ 0\rangle 3\rangle \text{swap} \rangle$	→	$ 0\rangle 3\rangle \text{swap} \rangle$
	→	$ 3\rangle 1\rangle$	→	$ 1\rangle 3\rangle \text{swap} \rangle$	→	$ 1\rangle 3\rangle \text{swap} \rangle$
	→	$ 3\rangle 2\rangle$	→	$ 2\rangle 3\rangle \text{swap} \rangle$	→	$ 2\rangle 3\rangle \text{swap} \rangle$
	→	$ 3\rangle 3\rangle$	→	$ 3\rangle 3\rangle \text{no swap} \rangle$	→	$3\rangle 3\rangle \text{no swap} \rangle$

Outcome:

$$\left(|0\rangle |1\rangle + |0\rangle |2\rangle + |0\rangle |3\rangle + |1\rangle |2\rangle + |1\rangle |3\rangle + |2\rangle |3\rangle \right) \otimes \left(| \text{swap} \rangle + | \text{no swap} \rangle \right)$$

Antisymmetrization via a Quantum Sort

Discard the sorted data register and use the register with swaps to (anti)symmetrize a given state

$(|4\rangle|7\rangle)(|\text{swap}\rangle + |\text{no swap}\rangle) \rightarrow (|4\rangle|7\rangle + |7\rangle|4\rangle)|0\rangle$
 by running the sort backwards. The comparisons in the sort will clean the ancillae encoding swaps.

$\rightarrow 0\rangle 0\rangle$	$\rightarrow 0\rangle 0\rangle \text{no swap}\rangle$	$\rightarrow \cancel{ 0\rangle 0\rangle \text{no swap}\rangle}$
$\rightarrow 0\rangle 1\rangle$	$\rightarrow 0\rangle 1\rangle \text{no swap}\rangle$	$\rightarrow \cancel{ 0\rangle 1\rangle \text{no swap}\rangle}$
$\rightarrow 0\rangle 2\rangle$	$\rightarrow 0\rangle 2\rangle \text{no swap}\rangle$	$\rightarrow 0\rangle 2\rangle \text{no swap}\rangle$
$\rightarrow 0\rangle 3\rangle$	$\rightarrow 0\rangle 3\rangle \text{no swap}\rangle$	$\rightarrow 0\rangle 3\rangle \text{no swap}\rangle$
$\rightarrow 1\rangle 0\rangle$	$\rightarrow 0\rangle 1\rangle \text{swap}\rangle$	$\rightarrow 0\rangle 1\rangle \text{swap}\rangle$
$\rightarrow 1\rangle 1\rangle$	$\rightarrow 1\rangle 1\rangle \text{no swap}\rangle$	$\rightarrow \cancel{ 1\rangle 1\rangle \text{no swap}\rangle}$
$\rightarrow 1\rangle 2\rangle$	$\rightarrow 1\rangle 2\rangle \text{no swap}\rangle$	$\rightarrow 1\rangle 2\rangle \text{no swap}\rangle$
$\rightarrow 1\rangle 3\rangle$	$\rightarrow 1\rangle 3\rangle \text{no swap}\rangle$	$\rightarrow 1\rangle 3\rangle \text{no swap}\rangle$
$\rightarrow 2\rangle 0\rangle$	$\rightarrow 0\rangle 2\rangle \text{swap}\rangle$	$\rightarrow 0\rangle 2\rangle \text{swap}\rangle$
$\rightarrow 2\rangle 1\rangle$	$\rightarrow 1\rangle 2\rangle \text{swap}\rangle$	$\rightarrow 1\rangle 2\rangle \text{swap}\rangle$
$\rightarrow 2\rangle 2\rangle$	$\rightarrow 2\rangle 2\rangle \text{no swap}\rangle$	$\rightarrow \cancel{ 2\rangle 2\rangle \text{no swap}\rangle}$
$\rightarrow 2\rangle 3\rangle$	$\rightarrow 2\rangle 3\rangle \text{no swap}\rangle$	$\rightarrow 2\rangle 3\rangle \text{swap}\rangle$
$\rightarrow 3\rangle 0\rangle$	$\rightarrow 0\rangle 3\rangle \text{swap}\rangle$	$\rightarrow 0\rangle 3\rangle \text{swap}\rangle$
$\rightarrow 3\rangle 1\rangle$	$\rightarrow 1\rangle 3\rangle \text{swap}\rangle$	$\rightarrow 1\rangle 3\rangle \text{swap}\rangle$
$\rightarrow 3\rangle 2\rangle$	$\rightarrow 2\rangle 3\rangle \text{swap}\rangle$	$\rightarrow 2\rangle 3\rangle \text{swap}\rangle$
$\rightarrow 3\rangle 3\rangle$	$\rightarrow 3\rangle 3\rangle \text{no swap}\rangle$	$\rightarrow \cancel{ 3\rangle 3\rangle \text{no swap}\rangle}$

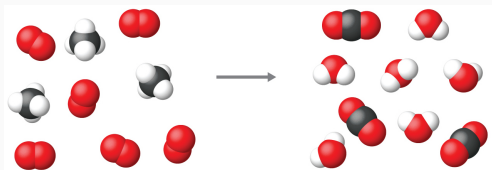
Complexity:

- $\mathcal{O}(m \log^2 m \log N)$ gates (prior: $\mathcal{O}(m^2 \log^2 N)$)
- $\mathcal{O}(\log^2 m \log \log N)$ depth (prior: $\mathcal{O}(m^2 \log^2 N)$)
- polynomial and exponential improvement to Abrams & Lloyd

Simulating time-dependent Hamiltonians

Motivation: Simulating Quantum Systems is Notoriously Difficult

First application of quantum computing



Product formulae: [Lloyd 1996, Aharonov & Ta-Shma 2003, Berry et al. 2007, Wiebe 2010, Campbell 2018, ...]

LCU: [Childs & Wiebe 2012, Berry et al. 2013, Berry et al. 2014, Berry et al. 2015, Berry et al. 2017, ...]

Other: [Childs 2010, Berry & Childs 2009, Low & Chuang 2016, Low & Chuang 2017, Low & Wiebe 2018, Gilyén 2019, ...]

+ simulations of special Hamiltonians, resource studies, applications

Time-dependent Hamiltonian Simulation

Implement the time evolution

$$i \frac{d}{dt} |\psi(t)\rangle = H(t) |\psi(t)\rangle$$

with logical error at most ϵ where H is a d -sparse Hermitian matrix.

[Lloyd 1996, Aharonov & Ta-Shma 2003, Wiebe et al. 2011, Poulin et al. 2011, Low & Wiebe 2018, Berry et al. 2019]

NEW: Algorithm for simulating time-dependent Hamiltonians with poly-logarithmic dependence on the inverse error.

Application: Interaction Picture Simulation

Reduced dependence on the norm in exchange for a time-dependent Hamiltonian

$$H = A + B$$

A – large norm, easy to implement, for example diagonal

B – small norm, must use regular Hamiltonian simulation methods

$$H_I(t) = e^{iAt} B e^{-iAt}$$

[Low, Wiebe 2019]

Discretization

- Split the evolution into r segments
- Truncated Dyson (Taylor) series

$$U(0, T/r) \approx \sum_{k=0}^K \frac{(-i)^k}{k!} \mathcal{T} \int_0^{T/r} dt H(t_k) \dots H(t_1)$$

- Discretize the integrals

$$U(0, T/r) \approx \sum_{k=0}^K \frac{(-iT/r)^k}{M^k k!} \sum_{j_1, \dots, j_k=0}^{M-1} \mathcal{T} H(t_{j_k}) \dots H(t_{j_1})$$

- Decompose Hamiltonians into linear combinations of unitaries

H_ℓ

$$H(t) = \gamma \sum_{\ell=0}^{L-1} H_\ell(t) \quad \text{or} \quad H(t) = \sum_{\ell=0}^{L-1} \alpha_\ell(t) H_\ell$$

LCU for Hamiltonian Simulation

The evolution can be approximated by a linear combination of unitaries

$$U(0, T/r) \approx \sum_j \beta_j V_j,$$

where β_j -s are the coefficients and V_j -s products of unitaries.

$U(0, T/r)$ is implemented through



and then boosted to almost perfect success rate using oblivious amplitude amplification.

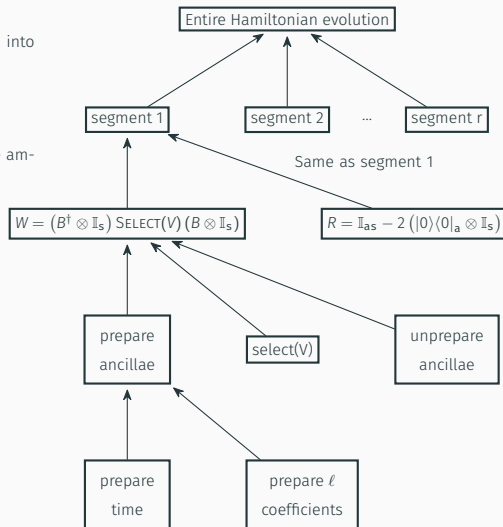
Overview

Divide the evolution into segments

Oblivious amplitude amplification

Circuits in the paper/thesis

State preparation



[Berry et al. 2014, Kieferova et al. 2019]

How Do We Enforce the Correct Order of Unitaries?

The times for the Hamiltonians must come in the correct order

$$\mathcal{T}H(t_{j_k}) \dots H(t_{j_2})H(t_{j_1}).$$

NEW: Two approaches to create a “clock” register.

Ordering based on sorting:

- value of k in unary
- create $\sum_i t_i \otimes \dots \otimes \sum_i t_i$
- sort the registers



Complexity:

- $\mathcal{O}\left(d^2 H_{\max} T \frac{\log(dH_{\max} T/\epsilon)}{\log \log(dH_{\max} T/\epsilon)}\right)$ queries, same as for the time-independent LCU algorithm

[Lloyd 1996, Aharonov 2003, Aharonov 2008, Wiebe et al. 2011, Poulin 2011] $\sim \text{poly}(1/\epsilon)$

- gate complexity depends on $\log(\|\dot{H}\|)$

([Lloyd 1996, Aharonov 2003, Aharonov 2008, Wiebe et al. 2011] polynomial dependence,

[Poulin 2011] independent)

Summary

- An efficient, low-depth algorithm for antisymmetrization of a fermionic wavefunction.
- A new algorithm simulating time-dependent Hamiltonians.

Thank you!